



Towards Context-based Epistemic Logic

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April 7, 2010

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Introduction

- **Purpose:**
- Context-based Epistemic Logic (CEL) = the logic of context-sensitive knowledge.
- (Traditional) Epistemic Logic (EL) = the logic of knowledge;
 - **CSK** is a notion abstracted from Epistemic Contextualism solution to Skepticism, which also being used to solve some other epistemic puzzles.

Introduction (cont'd)

- **Motivations:**

- **1.** To provide a new approach to the study of **limited rationality (LA)** in Epistemic Logic.
- Nowadays, the study of LA is extremely popular not only in the field of Epistemic Logic, but also in other fields like Game theory, Decision theory, Social software , and Artificial Intelligence (AI).

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- As far as I know, there are mainly two approaches to deal with LA in EL:
 - Through awareness: (Cf. Fagin, 1988; & de Jager, 2009);
 - Through accessible: (Cf. Pacuit et al, 09).



– 2. To strengthen further the connection between **epistemology** and **epistemic logic**.

– Quote:

“At first sight, the modern agenda of epistemology has little to do with logic... Now, epistemic logic started as a contribution to epistemology, or at least a tool in its modus operandi, with the seminal book *Knowledge and Belief* (Hintikka’s, 1962,2005).”

---from (van Benthem, 2006)

A scenic view of a beach with turquoise water, white sand, and a blue sky with white clouds and seagulls. The text "Skepticism and EC" is centered in the middle of the image.

Skepticism and EC

Skeptical Argument (SA)

- **Basic form** (DeRose, 1995):
 - P1: I don't know that not-H.
 - P2: If I don't know that not-H, then I don't know that O.
 - C: So, I don't know that O.

SA (cont'd)

- **Example:**
 - P1: I don't know that I am not a BIV.
 - P2: If I don't know that I am not a BIV, then I don't know that I have two hands.
 - C: So, I don't know that I have two hands.

Epistemic Contextualism (EC)

- **Quote:**

“...**EC** is the view that the proposition expressed by a given knowledge attribution (like “A knows that P” or “A knows that not P”) depends upon the context in which it is uttered.”

---from (Rysiew,2009).

EC's solution to Skepticism

- The presence of P1 has changed the context, such that a higher standard of knowledge are required.
- **Advantages:** explain the persuasiveness of SA & protect the correctness of our ordinary knowledge.

Further Question

- **Philosophers:** What is a context?
 - Cf. (Barke,2004); epistemic assumptions.
- **Logicians:** How to represent a context?
 - Cf. (Stalnaker,1998); a set of possible worlds (or states).



Epistemic Logic & Context Logic

Epistemic Logic (EL)

- **Language (L_E):**

$$\varphi ::= p \mid \neg\varphi \mid \varphi \wedge \psi \mid K\varphi;$$

Where $p \in P$.

Note: K is the abbreviation for K_a , since I only consider one agent a for simplicity.

EL (cont'd)

- **Epistemic model :**

$$M = \langle W, \approx, V \rangle$$

Where:

- W is a non-empty set;
- \approx is an equivalence relation on W ;
- V is a valuation mapping each $p \in P$ to a subset of W , i.e., $V(p) \in 2^W$.

EL (cont'd)

- **Semantics:**

$M, w \models p$, iff $w \in V(p)$;

$M, w \models K\phi$, iff for all $w' \in \approx_w$, $M, w' \models \phi$;

where $\approx_w = \{v \mid v \in W \& w \approx v\}$.

EL (cont'd)

- **Axiomatization (S5):**

- **Taut:** All instantiations of propositional tautologies;

- **K:** $K(\varphi \rightarrow \psi) \rightarrow (K\varphi \rightarrow K\psi)$;

- **T:** $K\varphi \rightarrow \varphi$; (Truth)

- **4:** $K\varphi \rightarrow KK\varphi$; (Positive introspection)

- **5:** $\neg K\varphi \rightarrow K\neg K\varphi$; (Negative introspection)

- **MP:** From φ and $\varphi \rightarrow \psi$, infer ψ ;

- **N:** From φ , infer $K\varphi$.

EL (cont'd)

- **Notation:**

- We denote EL-validity and EL-provability of φ as “ $\models_{\text{EL}} \varphi$ ” and “ $\vdash_{\text{EL}} \varphi$ ”, respectively.

- **Completeness of EL:**

- **Theorem 1:** For any φ , $\models_{\text{EL}} \varphi$ iff $\vdash_{\text{EL}} \varphi$.
- **Proof.** Cf. (van Ditmarsch et al, 2007, Chapter 7).

Context Logic (CL*)

- **Language (L_C):**

$$\varphi ::= p \mid \neg\varphi \mid \varphi \wedge \psi \mid [X]\varphi \mid [U]\varphi;$$

Where $p \in P$, and $X \in C$.

C is the index set of contexts, $[X]$ is context operator, $[U]$ is the universal modality.

Duals:

$$\langle X \rangle \varphi = \neg [X] \neg \varphi; \quad \langle U \rangle \varphi = \neg [U] \neg \varphi.$$

CL (cont'd)

- **Context model :**

$$M = \langle W, R, V \rangle$$

Where:

- W is a non-empty set;
- R is a function mapping each $X \in C$ to a subset of W , i.e., $R(X) \in 2^W$; ***Henceforth we write R_X for $R(X)$.**
- V is a valuation mapping each $p \in P$ to a subset of W , i.e., $V(p) \in 2^W$.

CL (cont'd)

- **Semantics:**

$M, w \models p$, iff $w \in V(p)$;

$M, w \models [X]\varphi$, iff for all $w' \in R_x$, $M, w' \models \varphi$;

$M, w \models [U]\varphi$, iff for all $w' \in W$, $M, w' \models \varphi$.

CL (cont'd)

- **Axiomatization (K45^{XY}):**
 - **Taut** plus the following, where $X, Y \in C \cup \{U\}$:
 - **K^X**: $[X](\phi \rightarrow \psi) \rightarrow ([X]\phi \rightarrow [X]\psi)$;
 - **T^U**: $[U]\phi \rightarrow \phi$;
 - **4^{XY}**: $[X]\phi \rightarrow [Y][X]\phi$;
 - **5^{XY}**: $\langle X \rangle \phi \rightarrow [Y]\langle X \rangle \phi$;
 - **MP**: From ϕ and $\phi \rightarrow \psi$, infer ψ ;
 - **N^X**: From ϕ , infer $[X]\phi$.

CL (cont'd)

- **Notation:**

- We denote CL-validity and CL-provability of φ as “ $\models_{\text{CL}}\varphi$ ” and “ $\vdash_{\text{CL}}\varphi$ ”, respectively.

- **Completeness of CL:**

- **Theorem 2:** For any φ , $\models_{\text{CL}}\varphi$ iff $\vdash_{\text{CL}}\varphi$.
- *Proof.* Cf. (Grossi et al, 2008).



Foundations of CEL

Foundations of CEL

- **Context-based epistemic model :**

$$M = \langle W, R, \approx, V \rangle$$

Where:

- W is a non-empty set;
- R is a function mapping each $X \in C$ to a subset of W , i.e., $R_X \in 2^W$;
- \approx is an equivalence relation on W ;
- V is a valuation mapping each $p \in P$ to a subset of W , i.e., $V(p) \in 2^W$.

Foundations of CEL (cont'd)

- **Definitions of CSK:**

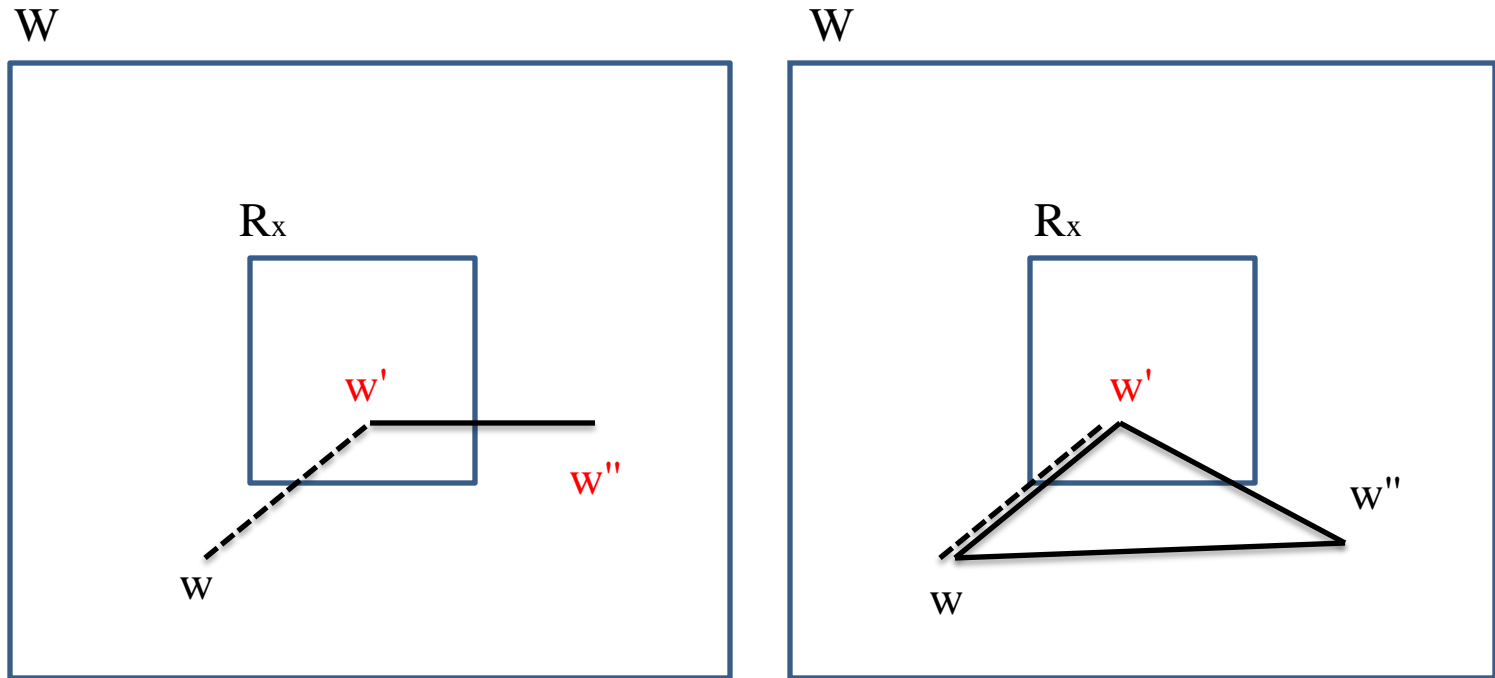
- **Static style:**

- $M, w \models [X]K\phi$, iff for all $w' \in R_X$, $M, w' \models K\phi$,
iff for all $w'' \in \approx_w$, $M, w'' \models \phi$.

- **Dynamic style:**

- $M, w \models K^{[X]}\phi$, iff for all $w' \in \approx_w \cap R_X$, $M, w' \models \phi$.
 - If we allow X to be U , then the standard epistemic operator K revives as $K^{[U]}$.

Contrasts



Static: $w \models [X]K\phi$, iff $w' \models \phi$ and $w'' \models \phi$.

Dynamic: $w \models K^{[X]}\phi$, iff $w' \models \phi$.

Figure 1. Definitions of CSK: Static vs. Dynamic

Contrasts (cont'd)

- **For static-style definition:**
 - (i) CSK is defined in term of context-free knowledge;* (which seems to be questionable.)
 - (ii) $K\phi \rightarrow [X]K\phi$ is not valid;
 - (iii) Nonetheless, $[X]K\phi \rightarrow [X]\phi$ is valid.
- **For dynamic-style definition:**
 - (i) CSK is defined independently;
 - (ii) $K\phi \rightarrow K^{[X]}\phi$ is valid;
 - (iii) Nonetheless, $K^{[X]}\phi \rightarrow [X]\phi$ is not valid.

Contrasts (cont'd)

- **Static-style** seems to be consistent with an **objective** understanding of context, since $[X]K\phi \rightarrow [X]\phi$ is valid. (*However, I haven't discovered any concrete example yet.)
- **Dynamic-style** seems to work well with the **subjective** understanding of context (esp. as **common assumptions**). Example: (Cf. next page)

SA in Dynamic CSK

- **For example:**

p: I am not a BIV;

q: I feel that I have two hands;

r: I have to hands.

The model is indicated as **Figure 2**
on the right side, where

$V(p)=\{u, v\}$; $V(q)=\{u, s\}$; $V(r)=\{u\}$.

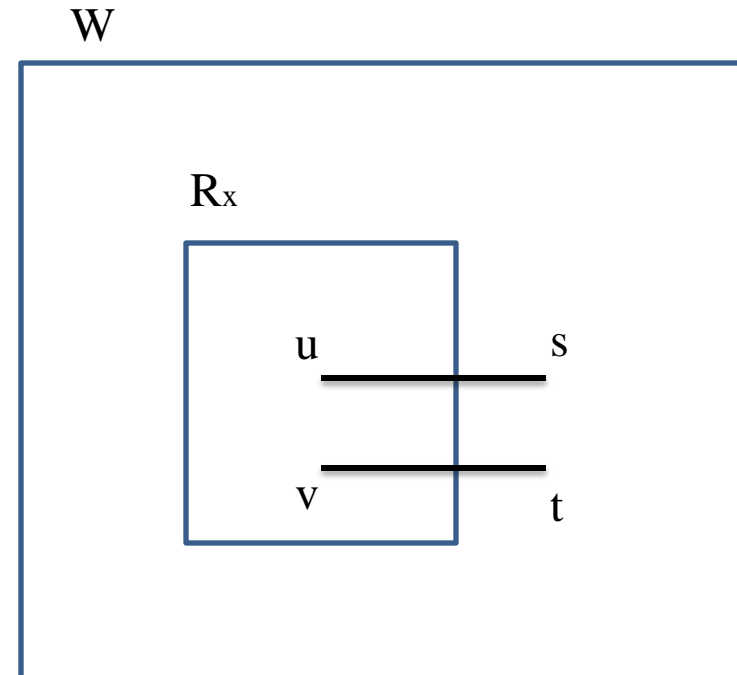


Figure 2. SA in Dynamic CSK

SA in Dynamic CSK (cont'd)

- (i) In all states, I know my feelings (i.e., either I feel that I have two hands or not).
- (ii) If w is the real state, then with the assumption of X , I know that I have two hands and I am not a BIV.
- (iii) When $P1$ of SA appears, the context has been extended to contain either s or t (depending upon which of u and v take as the actual state).
- (iv) So, the extension of context corresponds to the retraction of assumption. *(In next paper, I will revisit this example with more details after I work out the whole dynamic mechanics of DCEL.)



Candidates: SCEL & DCEL

SCEL

- **Language (L_{SCE}):**

$$\varphi ::= p \mid \neg\varphi \mid \varphi \wedge \psi \mid K\varphi \mid [X]\varphi \mid [U]\varphi;$$

Where $p \in P$, and $X \in C$.

- **Semantics:**

$M, w \models K\varphi$, iff for all $w' \in \approx_w$, $M, w' \models \varphi$;

$M, w \models [X]\varphi$, iff for all $w' \in R_X$, $M, w' \models \varphi$;

$M, w \models [U]\varphi$, iff for all $w' \in W$, $M, w' \models \varphi$.

SCEL (cont'd)

- **Axiomatization:**

- All axioms and rules of EL and CL, plus the axiom schemas below:

- 4^{XK} : $[X]\varphi \rightarrow K[X]\varphi$;

- 5^{XK} : $\langle X \rangle \varphi \rightarrow K \langle X \rangle \varphi$.

***Remark:** Knowledge operator is semi-context, since the following schemas are generally invalid:

- 4^{KY} : $K\varphi \rightarrow [Y]K\varphi$;

- 5^{KY} : $\langle K \rangle \varphi \rightarrow [Y]\langle K \rangle \varphi$.

SCEL (cont'd)

- **Notation:**

- We denote SCEL-validity and SCEL-provability of φ as “ $\models_{\text{SCEL}}\varphi$ ” and “ $\vdash_{\text{SCEL}}\varphi$ ”, respectively.

- **Completeness of CL:**

- **Theorem 3:** For any φ , $\models_{\text{SCEL}}\varphi$ iff $\vdash_{\text{SCEL}}\varphi$.
- *Proof.* Cf. (Xu, forthcoming).

DCEL

- **Language (L_{DCE}):**

$$\varphi ::= p \mid \neg\varphi \mid \varphi \wedge \psi \mid [X]\varphi \mid [U]\varphi \mid K^{[X]}\varphi \mid K^{[U]}\varphi;$$

Where $p \in P$, and $X \in C$.

- **Semantics:**

$M, w \models [X]\varphi$, iff for all $w' \in R_X$, $M, w' \models \varphi$;

$M, w \models [U]\varphi$, iff for all $w' \in W$, $M, w' \models \varphi$;

$M, w \models K^{[X]}\varphi$, iff for all $w' \in \approx_w \cap R_X$, $M, w' \models \varphi$;

$M, w \models K^{[U]}\varphi$, iff for all $w' \in \approx_w \cap W$, $M, w' \models \varphi$.

Notice that, $K^{[U]}$ is the same as K of SCEL, so DCEL is an extension of SCEL, and henceforth we write $K^{[U]}$ as K .

DCEL (cont'd)

- **Axiomatization:**

- All axioms and rules of SCEL, plus the axiom schemas and rule below, where $X \in C$:

- $\mathbf{K}^{\mathbf{K}[X]}$: $\mathbf{K}^{[X]}(\varphi \rightarrow \psi) \rightarrow (\mathbf{K}^{[X]}\varphi \rightarrow \mathbf{K}^{[X]}\psi)$;

- $\mathbf{K}\varphi \vee [X]\varphi \rightarrow \mathbf{K}^{[X]}\varphi$;

- $\mathbf{N}^{\mathbf{K}[X]}$: From φ , infer $\mathbf{K}^{[X]}\varphi$.

DCEL (cont'd)

- **Notation:**
 - We denote DCEL-validity and DCEL-provability of φ as “ $\models_{\text{DCEL}} \varphi$ ” and “ $\vdash_{\text{DCEL}} \varphi$ ”, respectively.
- **Completeness of DCEL:**
 - **Theorem 3:** For any φ , $\models_{\text{DCEL}} \varphi$ iff $\vdash_{\text{DCEL}} \varphi$.
 - *Proof.* Cf. (Xu, forthcoming).

Conclusion

- I have introduced the philosophical background of CSK and preliminaries of CEL: EL & CL;
- After that, I have proposed two distinct ways of defining CSK and made some detailed contrast;
- Further, I have obtained two candidate systems of CEL (namely, SCEL & DCEL) and proved their completeness.

Future Work

- **Future work:**
 - Develop the dynamic version of SCEL and DCEL;
 - (E.g., DCL; Cf. Aucher et al, 2009).
 - Compare with update semantics;
 - (Cf., Veltman, 1996; & de Jager, 2009).
 - Extend with more philosophical discussion;
 - (Cf. Lewis, 1996).
 - ...

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A scenic view of a beach with turquoise water, white sand, and a blue sky with white clouds and seagulls. The text "Thank you!" is overlaid in the center in a bold, yellow, serif font.

Thank you!